

Formal Connected Basic Pairs

Marco Benini

Dipartimento di Informatica e Comunicazione
Università degli Studi dell'Insubria

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What is a connected space?

- Following J.R. Munkres, *Topology*, Prentice Hall (2000, 2nd edition):

One says that a space can be “separated” if it can be broken up into two “globs” — disjoint open sets.

Otherwise, one says that it is connected.

Definition (Connected space)

Let X be a topological space.

A *separation* of X is a pair A, B of disjoint nonempty open subsets of X whose union is X .

The space X is said to be *connected* if there does not exist a separation of X .

A few comments

- Definition I is *negative*: a connected set does **not** have the property of being separated.
- Definition I implies a quantification over open subsets. And we cannot limit the quantification to elements of a basis.
- Definition I requires points. In fact, A and B must be nonempty (having at least one point) and disjoint (no points in common).
- The goal of this talk is to give a point-free definition of connected space and connected (sub)set in the framework of formal topology, specifically in Sambin's Basic Picture.

Connected subsets

Definition (Connected subset)

A subset $E \subseteq X$ is said to be *connected* in the topological space X if the corresponding subspace E is connected.

- The definition is easy and natural
- It requires the auxiliary notion of subspace.

Formal connected spaces

Definition (Connected space)

Let X be a topological space. A *separation* of X is a pair A, B of disjoint nonempty open subsets of X whose union is X .

The space X is said to be *connected* if there does not exist a separation of X .

- The first step in building a constructive version is to formulate Definition 1 in the language of basic pairs.
- We notice that, in classical logic, being A and B nonempty open subsets, the definition is equivalent to

$$X \subseteq A \cup B \rightarrow \exists x \in X. x \in A \cap B .$$

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$$X \subseteq A \cup B \rightarrow \exists x \in X. x \in A \cap B .$$

- Let's suppose to have the basic pair $\langle X; \Vdash; S \rangle$.
- A is an open subset if and only if $A = \text{ext} U$ for some $U \subseteq S$. Similarly, $B = \text{ext} V$.
- Moreover, supposing $S = \diamond X$, i.e., $\forall a \in S. \text{ext} a \not\ll X$, the open subset $A = \text{ext} U$ is nonempty if and only if U is nonempty. Similarly for V .
- The previous condition can be rephrased as, given U and V nonempty subsets of S ,

$$X \subseteq \text{ext} U \cup \text{ext} V \rightarrow \text{ext} U \not\ll \text{ext} V .$$

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$$X \subseteq \text{ext}U \cup \text{ext}V \rightarrow \text{ext}U \checkmark \text{ext}V .$$

- Thus, supposing that $X = \text{ext}S$, for nonempty U and V , the condition is equivalent to

$$\text{ext}S \subseteq \text{ext}(U \cup V) \rightarrow \text{ext}U \checkmark \text{ext}V .$$

- By definition of \square , the condition can be rephrased as

$$S \subseteq (\square \circ \text{ext})(U \cup V) \rightarrow \text{ext}U \checkmark \text{ext}V .$$

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$$S \subseteq (\square \circ \text{ext})(U \cup V) \rightarrow \text{ext}U \checkmark \text{ext}V .$$

- By definition of \mathcal{A} and \diamond ,

$$S \subseteq \mathcal{A}(U \cup V) \rightarrow U \checkmark (\diamond \circ \text{ext})V .$$

- It becomes natural to define a new operator $\text{conn} = \diamond \circ \text{ext}$:

$$S \subseteq \mathcal{A}(U \cup V) \rightarrow U \checkmark \text{conn}V .$$

- It is immediate to show that the previous condition is equivalent to the symmetrical

$$S \subseteq \mathcal{A}(U \cup V) \rightarrow V \checkmark \text{conn}U .$$

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- The previous condition encodes the notion of connected space if $X = \text{ext}S$, $S = \diamond X$ and both U and V are nonempty.
- The first and the second assumptions are *structural* so it makes sense to single out basic pairs meeting them:

Definition (Spatial basic pair)

A basic pair $\langle X; \Vdash; S \rangle$ is *spatial* if $X = \text{ext}S$ and $S = \diamond X$.

- A spatial basic pair represents a space having no point outside any open subset (since $X = \text{ext}S$), and no empty basic neighbourhood (since $S = \diamond X$).

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- Intuitively, $\text{conn}A$ is the collection of basic neighbourhoods having a point in common with A , that is, $a \in \text{conn}A$ if and only if a is *directly connected* to A .
- If, by symmetry, we define $\text{nocc} = \text{ext} \circ \diamond$, it is possible to describe in a slightly different way the condition of being spatial.
- In fact, it is easy to prove that the assumptions $S = \diamond X$ and $X = \text{ext}S$ are equivalent to $S = \text{conn}S$ and $X = \text{nocc}X$, respectively.
- Although this alternative definition has an appeal, its content is the same as the given one and no benefit is gained adopting it.

Formal connected spaces

Definition (Formal connected space)

A spatial basic pair $\langle X; \Vdash; S \rangle$ is *connected* if and only if the following condition holds for any U, V subsets of S :

$$U \neq \emptyset \wedge V \neq \emptyset \wedge S = \mathcal{A}(U \cup V) \rightarrow U \checkmark \text{conn} V .$$

- The condition is *formal*, i.e., it lives only in the formal side of the basic pair.
- It is *reasonable* since the condition says that, when U and V are nonempty and their union covers S , then the collection of basic neighbourhoods directly connected to V has an element in common with U .

Subspaces

- Spatial basic pairs have a natural definition that closely fits our intuition.
- But they also suggest how to define a coherent notion of subspace.

Definition (Subspace)

Given a spatial basic pair $\mathbb{S} = \langle X; \Vdash; \mathcal{S} \rangle$ and a subset $Y \subseteq X$, we say that the subspace on Y in \mathbb{S} is the basic pair $\mathbb{Y} = \langle X; \Vdash_Y; \mathcal{S} \rangle$ where $x \Vdash_Y a \equiv x \Vdash a \wedge x \in Y$.

- Evidently, $\text{ext}_Y U = \text{ext} U \cap Y$ so the open subsets of \mathbb{Y} are exactly the open sets in the classical construction of the subspace topology over Y .

Subspaces

- A few calculations lead to prove that $\text{conn}_Y S = \diamond Y$, that is, the basic neighbourhoods immediately connected to S in \mathbb{Y} are the basic neighbourhoods having at least one point in Y (in \mathbb{S}).
- Similarly, one proves that $\text{nocc}_Y X = \text{ext} S \cap Y$, that is, the points of X which are “glued” to some other point via the basic neighbourhoods of \mathbb{Y} are exactly the points of Y lying in some basic neighbourhood of S in \mathbb{S} .
- Since \mathbb{S} is spatial, $\text{ext} S \cap Y = X \cap Y = Y$, so $\text{nocc}_Y X = Y$.
- Therefore, the “domain” of the \Vdash_Y relation is exactly $\text{nocc}_Y X$; analogously, the domain of \Vdash is $\text{nocc} X$.
- Similarly, the “range” of the \Vdash_Y relation is exactly $\text{conn}_Y S$; analogously, the range of \Vdash is $\text{conn} S$.

Subspaces

- Thus, a subspace $\langle X; \Vdash_Y; S \rangle$ is thought as the “topology” on the points corresponding to the domain of \Vdash_Y , whose basis is composed by the extensions of the basic neighbourhoods in the range of \Vdash_Y .
- Since the domain of \Vdash_Y is $\text{nocc}_Y X = \text{ext} S \cap Y = Y$ and the range of \Vdash_Y is $\text{conn}_Y S = \diamond Y$, we are able to express the relevant notions of our interpretation of the subspace both inside its language and in the language of the space.

Formal connected subspaces

- To define the notion of connected subspace, we want to say that, for every pair of nonempty open subsets in the subspace, if their union is the domain of $\|_{-Y}$, then their intersection must contain at least one point.
- Since we can prove that $\text{ext}_Y U = Y \cap \text{ext} U$ and $\text{ext} U \overset{\circ}{\cap} Y \equiv U \overset{\circ}{\cap} \text{conn}_Y S$, which is a condition lying on the formal side of the subspace, we can define:

Definition (Formal connected subspace)

A subspace $\langle X; \|_{-Y}; S \rangle$ is *connected* if and only if the following condition holds for any U, V subsets of S :

$$U \overset{\circ}{\cap} \text{conn}_Y S \wedge V \overset{\circ}{\cap} \text{conn}_Y S \wedge S = \mathcal{A}_Y(U \cup V) \rightarrow U \overset{\circ}{\cap} \text{conn}_Y V .$$

Formal connected spaces

- The previous definition suggests to abstract and to generalise:

Definition (Formal connected basic pair)

A basic pair $\langle X; \Vdash; S \rangle$ is *connected* if and only if the following condition holds for any U, V subsets of S :

$$U \not\ll \text{conn} S \wedge V \not\ll \text{conn} S \wedge S = \mathcal{A}(U \cup V) \rightarrow U \not\ll \text{conn} V .$$

- Evidently, this definition applies to general basic pairs, not just spatial basic pairs or subspaces.

Formal connected spaces

Definition (Formal connected basic pair)

A basic pair $\langle X; \Vdash; S \rangle$ is *connected* if and only if the following condition holds for any U, V subsets of S :

$$U \not\ll \text{conn} S \wedge V \not\ll \text{conn} S \wedge S = \mathcal{A}(U \cup V) \rightarrow U \not\ll \text{conn} V .$$

- **(interpretation)** We stipulate that the “points” are the elements in the domain of \Vdash , while the basis of the “topology” is denoted by the range of \Vdash . Hence the collection of points is $\text{ext} S$ while the basis is $\diamond X$.

Formal connected spaces

Definition (Formal connected basic pair)

A basic pair $\langle X; \Vdash; \mathcal{S} \rangle$ is *connected* if and only if the following condition holds for any U, V subsets of \mathcal{S} :

$$U \not\ll \text{conn} \mathcal{S} \wedge V \not\ll \text{conn} \mathcal{S} \wedge \mathcal{S} = \mathcal{A}(U \cup V) \rightarrow U \not\ll \text{conn} \mathcal{S} .$$

- Since $U \not\ll \text{conn} \mathcal{S} \equiv \text{ext} U \not\ll \text{ext} \mathcal{S}$, U denotes an open subset having at least one point, i.e., it is “nonempty”.
- Since $\mathcal{S} = \mathcal{A}(U \cup V) \equiv \text{ext} \mathcal{S} = \text{ext} U \cup \text{ext} V$, the union of the points in U and V is the collection of all points.
- Finally, $U \not\ll \text{conn} \mathcal{S} \equiv \text{ext} U \not\ll \text{ext} \mathcal{S}$, so there is a point lying in the open subsets denoted by U and V .

Conclusions

- We have shown a point-free definition of connected space and connected subset valid for generic basic pairs.
- The definition, under the correct interpretation, is natural and intuitive (well, if one knows the Basic Picture).
- As a “subproduct”, we defined subspaces as a special construction over basic pairs.
- Most of the elementary results about connected spaces can be easily proved using the given definition.

Future directions

- Path-connection is easy if we define paths as functions.
- But the Basic Picture provides a clean and elegant notion of continuous relation...
- ... which naturally suggests to use “relational paths”.
- We are exploring relational paths and the topic is NOT immediate.
- But this is for another talk, sometime in the future...

Thanks

- First and most important to **Giovanni Sambin** since, without the Basic Picture, nothing in this talk would make sense.
- **Brunella Gerla**, my colleague who shares the efforts behind this work and owes its merits.
- **Peter Schuster**, who kindly invited me in this workshop and who constantly reviews my ideas (and my mistakes).

Lemma

$$\text{conn}S = \diamond X.$$

Proof.

$$a \in \text{conn}S \equiv a \in (\diamond \circ \text{ext})S \equiv \exists x \in X. a \in \diamond x \wedge x \in \text{ext}S \equiv \\ \exists x \in X. a \in \diamond x \wedge \diamond x \checkmark S.$$

$$\text{But } a \in S, \text{ so } a \in \text{conn}S \equiv \exists x \in X. a \in \diamond x \equiv \exists x \in X. x \in \text{ext}a \equiv \\ \text{ext}a \checkmark X \equiv a \in \diamond X. \quad \square$$

Lemma

$$\text{noccS} = \text{extS}.$$

Proof.

$$x \in \text{noccX} \equiv x \in (\text{ext} \circ \diamond)X \equiv \exists a \in \mathbf{S}. x \in \text{ext}a \wedge a \in \diamond\mathbf{S} \equiv$$

$$\exists a \in \mathbf{S}. x \in \text{ext}a \wedge \text{ext}a \notin X.$$

But $x \in X$, so $x \in \text{noccX} \equiv \exists a \in \mathbf{S}. x \in \text{ext}a \equiv \exists a \in \mathbf{S}. a \in \diamond x \equiv$
 $\diamond x \notin \mathbf{S} \equiv x \in \text{extS}. \quad \square$

Slide 18 and 21

Lemma

$$\text{ext}_Y U = Y \cap \text{ext} U.$$

Proof.

$$\begin{aligned} x \in \text{ext}_Y U &\equiv \exists a \in U. x \in \text{ext}_Y a \equiv \exists a \in U. x \Vdash_Y a \equiv \exists a \in U. x \in \\ &Y \wedge x \Vdash a \equiv x \in Y \wedge \exists a \in U. x \Vdash a \equiv x \in Y \wedge x \in \text{ext} U \equiv x \in \\ &Y \cap \text{ext} U. \quad \square \end{aligned}$$

Lemma

$$\text{conn}_Y S = \diamond Y.$$

Proof.

$$a \in \text{conn}_Y S \equiv \exists x \in X. x \Vdash a \wedge x \in Y \wedge x \in \text{ext}_Y S \equiv \exists x \in Y. x \in \text{ext} a \wedge x \in \text{ext} S.$$

But $a \in S$, so

$$a \in \text{conn}_Y S \equiv \exists x \in Y. x \in \text{ext} a \equiv \text{ext} a \Vdash Y \equiv a \in \diamond Y. \quad \square$$

Lemma

$$\text{nocc}_Y X = Y \cap \text{ext} S.$$

Proof.

$$x \in \text{nocc}_Y X \equiv \exists a \in S. x \Vdash a \wedge x \in Y \wedge a \in \diamond_Y X \equiv \exists a \in S. x \Vdash a \wedge x \in Y \wedge (\exists z \in X. z \Vdash a \wedge z \in X).$$

$$\text{But } x \in X, \text{ so } x \in \text{nocc}_Y X \equiv \exists a \in S. x \Vdash a \wedge x \in Y \equiv x \in Y \wedge x \in \text{ext} S \wedge x \in Y \cap \text{ext} S. \quad \square$$

Lemma

$$U \checkmark \text{conn}_Y S \equiv \text{ext} U \checkmark Y.$$

Proof.

$$U \checkmark \text{conn}_Y S \equiv \text{ext}_Y U \checkmark \text{ext}_Y S \equiv \exists x \in X. x \in \text{ext}_Y U \wedge x \in \text{ext}_Y S \equiv \exists x \in X. (\exists a \in U. x \Vdash a \wedge x \in Y) \wedge (\exists b \in S. x \Vdash b \wedge x \in Y).$$

$$\text{But } U \subseteq S, \text{ so } U \checkmark \text{conn}_Y S \equiv \exists x \in X, a \in U. x \Vdash a \wedge x \in Y \equiv \exists x \in X. x \in \text{ext} U \wedge x \in Y \equiv \text{ext} U \checkmark Y. \quad \square$$